

MATH 542 NUMBER THEORY  
 Problems to Think About #2  
 CH. 2, #1-3

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January 31st, 2014

**(1)**

$f, g$  multiplicative

$$h(n) = \sum_{d|P(n)} f(d)g\left(\frac{n}{d}\right)$$

$P(n) = r_1 r_2 \cdots r_v$  distinct primes that  $|n$

$$P(1) = 1$$

Let  $(a, b) = 1$

Let  $a = p_1^{\alpha_1} \cdots p_k^{\alpha_k}$   $p_i$  distinct,  $\alpha_i > 0$

$$b = q_1^{\beta_1} \cdots q_l^{\beta_l}$$

$q_j$  distinct,  $\beta_j > 0$

$$p_i \neq q_j \quad \forall i, j$$

Let  $\bar{a} = p_1 \cdots p_k, \bar{b} = q_1 \cdots q_l \quad \therefore (\bar{a}, \bar{b}) = 1$

First, we show  $d|\bar{a}\bar{b} \iff d = d' d''$  where  $d'|a, d''|b$

$$\begin{aligned} (\Rightarrow) d|\bar{a}\bar{b} \Rightarrow d = p_{i_1} \cdots p_{i_m} \cdot q_{j_1} \cdots q_{j_n} \\ \text{where } i_s \in \{1, 2, \dots, k\} \\ \text{where } j_t \in \{1, 2, \dots, l\} \end{aligned}$$

$$\text{Let } d' = p_{i_1} \cdots p_{i_m}$$

$$d'' = q_{j_1} \cdots q_{j_n}$$

$$\Rightarrow d = d' d''$$

and  $d'|\bar{a}, d''|\bar{b}$

$$(\Leftarrow) d = d' d'' \text{ where } d'|\bar{a}, d''|\bar{b}$$

$$\Rightarrow \bar{a} = d' k_1, \quad \bar{b} = d'' k_2 \quad \text{some } k_1, k_2 \in \mathbb{N}$$

$$\Rightarrow \bar{a}\bar{b} = d' d''(k_1 k_2) \Rightarrow d = d' d''|\bar{a}\bar{b} \quad \checkmark$$

If  $ab = 1$ , define  $\bar{a}\bar{b} = 1$

$$\text{Now, } h(ab) = \sum_{d|P(ab)} f(d)g\left(\frac{ab}{d}\right) = \sum_{d|\bar{a}\bar{b}} f(d)g\left(\frac{ab}{d}\right)$$

$$= \sum_{d' d''|\bar{a}\bar{b}} f(d')g\left(\frac{a}{d'}\right)g\left(\frac{b}{d''}\right)$$

$$= \sum_{d'|\bar{a}} \sum_{d''|\bar{b}} f(d')g\left(\frac{a}{d'}\right)g\left(\frac{b}{d''}\right) \quad (\text{since } (\bar{a}, \bar{b}) = 1 \Rightarrow (d', d'') = 1 \text{ and since } (a, b) = 1 \Rightarrow \left(\frac{a}{d'}, \frac{b}{d''}\right) = 1)$$

$$= \sum_{d'|\bar{a}} f(d')g\left(\frac{a}{d'}\right) \sum_{d''|\bar{b}} f(d'')g\left(\frac{b}{d''}\right)$$

$$= h(a)h(b) \quad \therefore h \text{ is multiplicative}$$

$$\therefore F_h(t) = \sum_{n=1}^{\infty} h(n)t^n = \left(\sum_{k=0}^{\infty} h(p_1^k)t_1^k\right)\left(\sum_{k=0}^{\infty} h(p_2^k)t_2^k\right) \cdots$$

$$\begin{aligned}
h(p_i^k) &= \begin{cases} \sum_{d|1} f(d)g(\frac{1}{d}) = 1, & \text{when } k = 0. \\ \sum_{d|p_i} f(d)g(\frac{p_i^k}{d}), & \text{when } k > 0 \end{cases} \\
&= \begin{cases} 1, & \text{when } k = 0. \\ g(p_i^k) + f(p_i)g(p_i^{k-1}), & \text{when } k > 0 \end{cases} \\
\Rightarrow \sum_{k=0}^{\infty} h(p_i^k)t_i^k &= (1 + (g(p_i) + f(p_i))t_i + (g(p_i^2) + f(p_i)g(p_i))t_i^2 + \dots) \\
&= (1 + g(p_i)t_i + g(p_i^2)t_i^2 + \dots) + f(p_i)t_i(1 + g(p_i)t_i + g(p_i^2)t_i^2 + \dots) \\
&= F_{i,g}(t_i) + f(p_i)t_i F_{i,g}(t_i) \\
&= F_{i,g}(t_i)(1 + f(p_i)t_i) \\
\Rightarrow F_h(t) &= \prod_{i=1}^{\infty} F_{i,g}(t_i)(1 + f(p_i)t_i) = F_g(t) \prod_{i=1}^{\infty} (1 + f(p_i)t_i) \\
\text{Moreover, by noting } 1 + f(p_i)t_i &= 1 + f(p_i)\mu^2(p_i)t_i + f(p_i^2)\mu^2(p_i^2)t_i^2 + \dots \\
&= F_{i,f\mu^2}(t_i), \text{ we can write } F_h(t) = F_g(t)F_{f\mu^2}(t) \\
&\quad (f\mu^2 \text{ multiplicative since it is a product of multiplicative functions})
\end{aligned}$$

## (2)

Want  $\sum_{d|P(n)} f(d)\sigma(\frac{n}{d}) = 1$

Referring to ex. (1), we want  $h(n) = 1$

$$\therefore h = u$$

Want  $g = \sigma$

$\therefore$  We get, from the result of ex. (1),

$$F_u(t) = F_h(t) = F_g(t) \prod_{i=1}^{\infty} (1 + f(p_i)t_i) = F_{\sigma}(t) \prod_{i=1}^{\infty} (1 + f(p_i)t_i)$$

We know  $F_u(t) = \frac{1}{(1-t_1)(1-t_2)\dots}$

We know  $F_{\sigma}(t) = \frac{1}{(1-t_1)(1-t_2)\dots(1-p_1t_1)(1-p_2t_2)\dots}$

$$\Rightarrow f(p_i) = -p_i \forall i \geq 1 \quad (\text{and, of course, } f(1) = 1)$$

$\Rightarrow f(n) = \mu(n)I(n)$  is a possible candidate (f is multiplicative since it is the product of multiplicative functions)

Let's try it out on an example:

$$\text{Let } n = p_1^2 p_2$$

$$\begin{aligned}
&\Rightarrow \mu(1)(1)\sigma(p_1^2 p_2) + \mu(p_1)(p_1)\sigma(p_1 p_2) + \mu(p_2)(p_2)\sigma(p_1^2) + \mu(p_1 p_2)(p_1 p_2)\sigma(p_1) = \\
&(1 + p_1 + p_2 + p_1^2 + p_1 p_2 + p_1^2 p_2) - p_1(1 + p_1 + p_2 + p_1 p_2) - p_2(1 + p_1 + p_1^2) + p_1 p_2(1 + p_1) \\
&= 1 \quad \checkmark
\end{aligned}$$

## (3)

$$f : \mathbb{Z}^+ \rightarrow \mathbb{R} \ni f(n) = \chi(n \text{ odd})$$

Let  $n = ab$  where  $(a, b) = 1$

Since  $(a, b) = 1$ , at most one of  $a, b$  is even

(case 1)  $a$  odd,  $b$  odd  $\Rightarrow ab$  odd

$$f(ab) = \chi(ab \text{ odd}) = 1 = 1 \cdot 1 = \chi(a \text{ odd})\chi(b \text{ odd}) = f(a)f(b)$$

(case 2)  $a$  odd,  $b$  even  $\Rightarrow ab$  even

$$f(ab) = \chi(ab \text{ odd}) = 0 = 1 \cdot 0 = \chi(a \text{ odd})\chi(b \text{ odd}) = f(a)f(b)$$

Similarly for  $a$  even,  $b$  odd

$\therefore f$  is multiplicative

First, let  $n$  odd  $\Rightarrow d$  odd  $\forall d|n$

$$\Rightarrow f(d) = 1$$

$$\Rightarrow h(n) = \sum_{d|n} f(d)\phi\left(\frac{n}{d}\right) = \sum_{d|n} \phi\left(\frac{n}{d}\right) = \sum_{d|n} \phi(d) = n \quad (\text{done in class and in text}) \quad \checkmark$$

Now, let  $n$  even. Let  $n = 2^l m$  where  $m$  odd and  $l > 0$

$$\text{We know } h(n) = (f * \phi)(n) = [t^n]F_f(t)F_\phi(t)$$

$$F_f(t) = \left( \sum_{k=0}^{\infty} f(p_1^k)t_1^k \right) \left( \sum_{k=0}^{\infty} f(p_2^k)t_2^k \right) \dots$$

$$= (1)(1 + t_2 + t_2^2 + \dots)(1 + t_3 + t_3^2 + \dots)$$

$$= \frac{1}{(1-t_2)(1-t_3)\dots}$$

$$\text{We have } F_\phi(t) = \frac{(1-t_1)(1-t_2)\dots}{(1-p_1t_1)(1-p_2t_2)\dots} \quad (\text{done in class})$$

$$\therefore F_f(t)F_\phi(t) = \frac{(1-t_1)}{(1-p_1t_1)(1-p_2t_2)\dots} = (1-t_1)F_I(t)$$

$$\therefore h(p_1^{e_1} p_2^{e_2} \dots) = [t_1^{e_1} t_2^{e_2} \dots](1-t_1)F_{1\ I}(t_1)F_{2\ I}(t_2)F_{3\ I}(t_3) \dots$$

$$= [t_1^{e_1} t_2^{e_2} \dots](1-t_1) \left( \sum_{k=0}^{\infty} p_1^k t_1^k \right) F_{2\ I}(t_2)F_{3\ I}(t_3) \dots$$

$$= [t_1^{e_1} t_2^{e_2} \dots] \left( \sum_{k=0}^{\infty} (p_1^k t_1^k - p_1^k t_1^{k+1}) \right) F_{2\ I}(t_2)F_{3\ I}(t_3) \dots$$

Now, let  $n = (2^l)(p_2^{e_2} p_3^{e_3} \dots) = 2^l m$  with  $l > 0$ ,  $m$  odd

What  $k$ 's produce a term involving  $t_1^l$ ?

Answer:  $k = l$  and  $k = l - 1$  making  $2^l t_1^l - 2^{l-1} t_1^l = 2^{l-1}(2-1)t_1^l = 2^{l-1}t_1^l$

$$\therefore 2^{l-1} = [t_1^l]((1-t_1)F_{1\ I}(t_1))$$

$$\therefore h(n) = [t_1^l]((1-t_1)F_{1\ I}(t_1))[t_2^{e_2} t_3^{e_3} \dots]F_{2\ I}(t_2)F_{3\ I}(t_3) \dots$$

$$= 2^{l-1}I(m) = 2^{l-1}m = \frac{1}{2}n \quad \checkmark$$