MATH 542 NUMBER THEORY Problems to Think About #5 CH. 5, #1-7

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(1)

Let us call:

$$\begin{aligned} \alpha : (a, b, c) \to (c, -b, a) \text{ for } f \in (1) \text{ or } (3) \\ \beta : (a, b, c) \to (a, b + 2ka, ak^2 + bk + c) \text{ for } f \in (2) \text{ or } (5) \text{ where } k = [[\frac{a-b}{2a}]] \\ \gamma : (a, b, c) \to (a, b - 2ka, ak^2 - bk + c) \text{ for } f \in (4) \text{ or } (7) \text{ where } k = [[\frac{a+b}{2a}]] \\ \delta : (a, b, c) \to (a, b + 2a, a + b + c) = (a, a, c) \text{ for } f \in (6) \end{aligned}$$

First, we need to show that each is unimodular: Recall $F = \begin{pmatrix} a & b/2 \\ b/2 & c \end{pmatrix}$ (α) Let $U_{\alpha} = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix} \Rightarrow \det U_{\alpha} = 1$ and $U_{\alpha}^{T}FU_{\alpha} = \begin{pmatrix} c & -b/2 \\ -b/2 & a \end{pmatrix}$ (β) Let $U_{\beta} = \begin{pmatrix} 1 & k \\ 0 & 1 \end{pmatrix} \Rightarrow \det U_{\beta} = 1$ and $U_{\beta}^{T}FU_{\beta} = \begin{pmatrix} a & 1/2(b+2ka) \\ 1/2(b+2ka) & ak^{2}+bk+c \end{pmatrix}$ (γ) Let $U_{\gamma} = \begin{pmatrix} 1 & -k \\ 0 & 1 \end{pmatrix} \Rightarrow \det U_{\gamma} = 1$ and $U_{\gamma}^{T}FU_{\gamma} = \begin{pmatrix} a & 1/2(b-2ka) \\ 1/2(b-2ka) & ak^{2}-bk+c \end{pmatrix}$ (δ) Let $U_{\delta} = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \Rightarrow \det U_{\delta} = 1$ and $U_{\delta}^{T}FU_{\delta} = \begin{pmatrix} a & 1/2(b-2ka) \\ 1/2(b-2ka) & ak^{2}-bk+c \end{pmatrix}$

 \therefore Each is unimodular.

Let the transformed BQF be $Ax^2 + Bxy + Cy^2$. $f \in (6) \ (a < c, b = -a) \xrightarrow{\delta} A < C, B = A \Rightarrow \text{reduced}.$ $f \in (3) \ (a = c, -a \le b < 0) \xrightarrow{\alpha} A = C, 0 < B \le A \Rightarrow \text{reduced}.$ $f \in (1) \ (a > c) \xrightarrow{\alpha} A < C \begin{cases} -A < B \le A \Rightarrow \text{reduced} \\ B = -A \xrightarrow{\delta} A < C, B = A \Rightarrow \text{reduced} \\ |B| > A \text{ (need to go to } \beta \text{ or } \gamma) \end{cases}$ In $(\beta), k = [[\frac{a-b}{2a}]] \Rightarrow \frac{a-b}{2a} - 1 \le k < \frac{a-b}{2a} \Rightarrow -a \le b + 2ka < a$ So, $a \to A, -A \le B < A$. In $(\gamma), k = [[\frac{a+b}{2a}]] \Rightarrow \frac{a+b}{2a} - 1 \le k < \frac{a+b}{2a} \Rightarrow -a < b - 2ka \le a$ So, $a \to A$, $-A < B \leq A$.

We will demonstrate $f \in (2)$ (a = c, b < -a) $(f \in (5)$ with (β) and $f \in (4)$ or (7) with (γ) are entirely analogous).

$$f \in (2) \ (a = c, b < -a) \xrightarrow{\beta} \begin{cases} A < C, -A < B < A \Rightarrow \text{reduced} \\ A < C, B = -A \xrightarrow{\delta} A < C, B = A \Rightarrow \text{reduced} \\ A = C, 0 \le B < A \Rightarrow \text{reduced} \\ A = C, -A \le B < 0 \xrightarrow{\alpha} A = C, 0 < B \le A \Rightarrow \text{reduced} \\ A = C, -A \le B < A \xrightarrow{\alpha} A < C \begin{cases} -A < B \le A \Rightarrow \text{reduced} \\ B = -A \xrightarrow{\delta} A < C, B = A \Rightarrow \text{reduced} \\ B = -A \xrightarrow{\delta} A < C, B = A \Rightarrow \text{reduced} \\ \text{otherwise, } |B| > A \text{ and we are back} \\ \text{to apply } \beta \text{ or } \gamma \end{cases}$$

Each time we get to a possible A > C (and then apply α), if not reduced we end up with A < C, |B| > A, but now A is strictly less than its predecessor. Also, $0 \le |B| \le A$ after each application of β or γ . Therefore, if f not already reduced, |B| will eventually equal 0 after a finite number of steps, in which case f is either reduced or reduced after one more application of α .

Therefore, f gets reduced after a finite number of steps in all cases.

(2)-(7) mathematica programs attached

(7)

 $p \equiv 1 \pmod{4}$

Just to summarize the theory of infinite descent for two squares, which is very similar to the case of four squares but simpler:

Since
$$\left(\frac{-1}{p}\right) = 1, \exists x \in [0, p-1) \ni x^2 \equiv -1 \pmod{p}.$$

 $\therefore \exists m \ni mp = x^2 + 1$. Furthermore, $m \in [1, p - 1]$ as in four square case. Find an m (by brute force).

The rest is identical to the case of four squares, but we do not have to worry about pairing up like parity addends when m is even, since m even implies both addends are odd or both are even.